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## SOME CONSEQUENCES OF A RESULT OF EHRENFEUCHT AND ROZENBERG (\*)

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Communiqué par J. BERSTEL

*Abstract.* — Recently, Ehrenfeucht and Rozenberg have proved that there are context-free languages that are not EDTOL languages. In the present note, some consequences of this result are pointed out with reference to certain open problems related to matrix languages.

*Résumé.* — Récemment, Ehrenfeucht et Rozenberg ont prouvé qu'il existe des langages indépendants du contexte qui ne sont pas des EDTOL-langages. Dans la présente note, nous indiquons quelques conséquences de ce résultat, en résolvant certains problèmes ouverts sur les langages matriciels.

### 1. NOTATIONS

In order to save space we shall omit the definitions and we shall specify only the notations we use:

—  $\mathcal{L}$  (EDTOL) is the family of languages generated by EDTOL systems (see [6]);

—  $\mathcal{M}_f$  is the family of languages generated by matrix grammars of finite index [1];

—  $\mathcal{SM}$  is the family of simple matrix languages [7] and  $\mathcal{SM}_f$  is the family of finite index languages in  $\mathcal{SM}$ ;

—  $\mathcal{L}_2$  is the family of context-free languages;

—  $\text{Ind}_m(L)$ ,  $\text{Ind}_{cf}(L)$  denote the index of a language  $L$  according to matrix grammars, respectively, to context-free grammars;

—  $D_i$  is the Dyck language over the vocabulary  $\{a_1, \dots, a_i, a'_1, \dots, a'_i\}$ , that is the language generated by the context-free grammar with the rules  $S \rightarrow SS$ ,  $S \rightarrow a_j S a'_j$ ,  $j = 1, 2, \dots, i$ ,  $S \rightarrow \lambda$  [16].

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## 2. OPEN PROBLEMS RELATED TO THE FAMILIES $\mathcal{M}_f, \mathcal{SM}$

The index of a matrix grammar was defined in [1] where it is proved that for any infinite language  $L$  in  $\mathcal{M}_f$ , the set of lengths of strings in  $L$  contains an infinite arithmetical progression (a similar result holds for languages in  $\mathcal{L}_2$ ). In [1] it is asked (problem  $A$ ) whether or not any context-free language has a finite index (clearly, according to matrix grammars).

The problem was partially solved by Salomaa [15] by proving that  $\text{Ind}_{cf}(D_1)$  is infinite. The existence of a context-free language of infinite index according to matrix grammars remained open.

In [3], Cremers, Mayer, Weiss claim that  $\text{Ind}_m(D_1)=2$ , but, as in [8] it was pointed out, the proof in [3] is wrong; the finiteness of  $\text{Ind}_m(D_1)$  remained open. Note that if any Dyck language would be in  $\mathcal{M}_f$ , then  $\mathcal{L}_2 \subset \mathcal{M}_f$  since the family  $\mathcal{M}_f$  is a full-AFL [11] and any language  $L$  in  $\mathcal{L}_2$  can be written as  $h(D_i \cap R)$ , where  $h$  is a homomorphism and  $R$  is a regular language [16].

In [7] Ibarra introduces the so-called simple matrix grammars, a class of grammars generating languages with many context-free-like properties. Similar properties (closure properties, decision results, pumping lemmas) were proved for the family  $\mathcal{M}_f$  too (see [11, 12, 14]). The relation between the families  $\mathcal{M}_f$  and  $\mathcal{SM}$  was formulated as an open problem in various contexts [9, 11, 12]; in fact, we obviously have  $\mathcal{M}_f - \mathcal{SM} = \emptyset$  since the language  $\{a^n b^n c^n \mid n \geq 1\}^*$  is in  $\mathcal{M}_f$  but not in  $\mathcal{SM}$  [7]. Therefore, the problem asks for a language in  $\mathcal{SM}$  which has an infinite matrix index. A stronger formulation [12] asks for context-free languages which are not in  $\mathcal{M}_f$  (problem related to the paper [3]).

In [9] we proved that  $\mathcal{SM}_f \subset \mathcal{M}_f$  and one asks whether or not  $\mathcal{SM}_f \subseteq \mathcal{SM}$  is a proper inclusion. The problem was solved in [10] by proving that  $D_1$  does not belong to  $\mathcal{SM}_f$ . In a similar way it follows that any  $D_i, i \geq 1$ , is not in  $\mathcal{SM}_f$ .

## 3. THE SETTLEMENT OF THE ABOVE OPEN PROBLEMS

The result in [4] can be easily used to solve all the above open problems. Indeed, we have:

**THEOREM:**  $\text{Ind}_m(D_i)$  is infinite for any  $i \geq 2$ .

*Proof:* In [2] one introduces the “parallel matrix grammars” as usual context-free matrix grammars with parallel derivation (each rule rewrites all occurrences of its left-hand side in the string to be rewritten). Let us denote by  $\mathcal{PM}$  the family of languages generated by these grammars. In [2] it is proved that  $\mathcal{M}_f \subset \mathcal{PM}$ , strict inclusion.

In [13] it is proved that, in fact, we have  $\mathcal{PM} = \mathcal{L}$  (EDTOL), therefore  $\mathcal{M}_f \subset \mathcal{L}$  (EDTOL). Following [4], there are context-free languages which are not in  $\mathcal{L}$  (EDTOL). Thus there are context-free languages which do not belong to  $\mathcal{M}_f$  as well. Consequently, any context-free generator (any context-free language whose smallest AFL containing it equals the family of context-free languages) is not a matrix language of finite index. Any language  $D_i$ ,  $i \geq 2$ , is a context-free generator (see example 5.1.1, p. 139 [5]), hence  $\text{Ind}_m(D_i)$  is infinite for any  $i \geq 2$ .

Consequently, (1)  $\mathcal{L}_2$  and  $\mathcal{M}_f$  are incomparable (clearly,  $\mathcal{M}_f$  contains non-context-free languages), (2)  $\mathcal{PM}$  and  $\mathcal{M}_f$  are incomparable, (3)  $\text{Ind}_m(D_i) = \text{Ind}_{cf}(D_i) = \infty$ ,  $i \geq 2$ , and thus all the problems listed in section 2 are settled.

It remains open the problem to find context-free languages  $L$  such that  $\text{Ind}_m(L) < \text{Ind}_{cf}(L)$  [or even  $\text{Ind}_m(L) < \text{Ind}_{cf}(L) = \infty$ ]. Particularly, we do not know whether  $\text{Ind}_m(D_1)$  is finite or not.

Some problems settled by the result in [4] were considered in [4] too. Perhaps there are other problems which can be solved using the same result.

*Note:* We are very indebted to the referee for pointing out a mistake in an earlier version of this paper.

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