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REMARKS ON DOL GROWTH SEQUENCES (*)

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Abstract. — *Two theorems are given characterizing the position of* DOL *and* PDOL *growth sequences among N-rational sequences.*

1. INTRODUCTION

A DOL system or a deterministic context-independent Lindenmayer system consists of an initial word ω and a set of productions $x \to \delta(x)$ which give for any letter x and thus also for any word a unique successor. The growth sequence of a DOL system is the sequence formed by the lengths of the words ω, $\delta(\omega)$, $\delta^2(\omega)$, ... DOL sequences have been investigated e. g. in Paz and Salomaa [6], Salomaa [8], Vitànyi [10], Ruohonen [7] and Karhumäki [4].

A sequence (r_n) is called N-rational if it can be represented in the form $r_n = P M^n Q$ where P is a row vector, M is a square matrix, Q is a column vector and the entries of P, M and Q are natural numbers. (The name N-rational comes from the general theory of rational series founded by M. P. Schützenberger.) Now it is easy to see that a DOL sequence is N-rational; in fact it has a representation $PM^n Q$ where Q consists merely of ones.

If a DOL sequence is not terminating, i. e. if $r_n \neq 0$ for every n, and if L is the largest of the lengths of the words $\delta(x)$ then obviously $r_{n+1}/r_n \leqq L$ for every n. If the system under consideration is such that $\delta(x)$ is always a non-empty word then this system is called a PDOL system and its growth sequence is called a PDOL sequence. Obviously a PDOL sequence is non-decreasing.

The goal of this paper is to illustrate the position of DOL and PDOL sequences among N-rational sequences. It will be seen that the satisfaction of an inequality $r_{n+1}/r_n \leqq L$ is characteristic for DOL sequences. Further it will be seen that it is not the non-negativity but the N-rationality of the sequence $(r_{n+1} - r_n)$ that makes a DOL sequence to be a PDOL sequence.

2. PRELIMINARIES

A DOL system is at triple $G = (X, \delta, \omega)$ where $X = \{ x_1, \ldots, x_k \}$ is an alphabet, δ is an endomorphism of the free monoid X^* and $\omega \in X^*$. The mapping δ is usually given by writing the oductions $x_i \to \delta(x_i)$ and the

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word ω is called the axiom. If $\delta(x_i) \neq \lambda$ for each i then G is called a PDOL system. The function

$$f_G(n) = \lg(\delta^n(\omega)),$$

where lg means word length is called the growth function of G.

A pair (X, δ) where X and δ are as above is called a DOL scheme.

Introducing the axiom vector

$$P = (\lg_1(\omega), \ldots, \lg_k(\omega))$$

and the growth matrix

$$M = \begin{pmatrix} \lg_1(\delta(x_1)) & \ldots & \lg_k(\delta(x_1)) \\ . & \ldots & . \\ \lg_1(\delta(x_k)) & \ldots & \lg_k(\delta(x_k)) \end{pmatrix}$$

of G (here $\lg_j$ denotes the number of letters x_j) we obtain

$$(\lg_1(\delta^n(\omega)), \ldots, \lg_k(\delta^n(\omega))) = PM^n$$

and

$$f_G(n) = PM^n(1, \ldots, 1)^T.$$

A sequence $(r(n))$ is called Z-rational (resp. N-rational) if

$$r(n) = PM^n Q = (p_1, \ldots, p_k)\begin{pmatrix} m_{11} & \ldots & m_{1k} \\ . & \ldots & . \\ m_{k1} & \ldots & m_{kk} \end{pmatrix}^n \begin{pmatrix} q_1 \\ . \\ q_k \end{pmatrix},$$

where all the entries are integers (resp. non-negative integers). If now G is a DOL system (resp. a PDOL system) then by the above $(f_G(n))$ is a special N-rational sequence called a DOL sequence (resp. a PDOL sequence).

It is known (Schützenberger [9]) that a sequence $(r(n))$ is Z-rational (N-rational) iff the series $\sum r(n) x^n$ is Z-rational (N-rational). If now $\sum r(n) x^n$ is a non-polynomial N-rational series then a theorem of Berstel [1] concerning its poles tells the following: there are a natural number p, algebraic numbers $A, A_1, \ldots, A_s$ ($A > 0$, $|A_j| < A$, $s \geqq 0$) and polynomials $H_0, \ldots, H_{p-1}, h_1, \ldots, h_s$ such that

$$r(i+np) = H_i(i+np)A^{i+np} + \sum_{j=1}^{s} h_j(i+np)A_j^{i+np}$$

for large values of n ($i = 0, \ldots, p-1$). In the case of a DOL sequence the polynomials H_i must have a common degree l because the quotients $r(n+1)/r(n)$ are bounded from above. We shall say that $(r(n))$ has the growth order $n^l A^n$.

3. THE PDOL SEQUENCES

LEMMA 1: *If $(f(n)) = (PM^n Q)$ is an N-rational sequence and P has positive entries then $(f(n))$ is a* DOL *sequence.*

Proof: Let $G = (\{x_1, \ldots, x_k\}, \delta, \omega)$ be a DOL system with axiom vector Q^T and growth matrix M^T. Define $G' = (\{x_1, \ldots, x_k, x\}, \delta', \omega')$ where

$$\delta'(x_i) = \delta(x_i)\, x^{\lg_1(\delta(x_i))(p_1-1)+\ldots+\lg_k(\delta(x_i))(p_k-1)},$$

$$\delta'(x) = \lambda,$$

$$\omega' = \omega\, x^{\lg_1(\omega)(p_1-1)+\ldots+\lg_k(\omega)(p_k-1)}.$$

Then obviously $f(n) = Q^T (M^T)^n P^T = f_{G'}(n)$.

THEOREM 1: *Let $(r(n))$ be an N-rational sequence. Then we can find natural numbers m and p and* DOL *sequences $(d_0(n)), \ldots, (d_{p-1}(n))$ such that*

$$r(m+i+np) = d_i(n) \qquad (i = 0, \ldots, p-1).$$

Proof: Let $r(n) = PM^n Q$ and let $G = (X, \delta, \omega)$ be a DOL system with axiom vector P and growth matrix M. Denote by X_n the set of letters occurring in $\delta^n(\omega)$. Then we can find numbers m and p such that $X_{m+i} = X_{m+i+np}$. We may of course suppose that no y_n is empty.

Introduce now the DOL systems

$$G_i = (X_{m+i}, \delta^p, \delta^{m+i}(\omega)) \qquad (i = 0, \ldots, p-1)$$

whose axiom vectors and growth matrices are denoted by P_i and M_i. Then obviously

$$r(m+i+np) = P_i M_i^n Q_i,$$

where Q_i is composed of those entries of Q corresponding to letters of X_{m+i}. By lemma 1 we may define $d_i(n) = P_i M^n Q_i$.

THEOREM 2: *The following conditions are equivalent for a sequence $(r(n))$:*

(i) *$(r(n))$ is a* PDOL *sequence not identically zero;*

(ii) *$r(0)$ is a positive integer and the sequence $(s(n)) = (r(n+1) - r(n))$ is N-rational.*

Proof: Suppose (i) holds. If now $(r(n))$ corresponds to a PDOL system $G = (\{x_1, \ldots, x_k\}, \delta, \omega)$ then

$$s(n) = \sum_{i=1}^{k} \lg_i(\delta^n(\omega))(\lg(\delta(x_i)) - 1)$$

and each of the sequences $(\lg_i(\delta^n(\omega)))$ $(i = 1, \ldots, k)$ is N-rational.

Suppose then that (ii) holds. Write according to theorem 1

$$s(m+i+np) = d_i(n) \qquad (i = 0, \ldots, p-1)$$

where $(d_i(n))$ corresponds to a system $G_i = (X_i, \delta_i, \omega_i)$. Assuming that the alphabets X_i are mutually disjoint we construct the PDOL system $G = (X, \delta, \omega)$ where

$$X = \left(\bigcup_{i=0}^{p-1} \bigcup_{j=0}^{p-1} X_i^{(j)}\right) \cup \{y\},$$

$$\omega = \omega_0^{(p-1)} \omega_1^{(p-2)} \ldots \omega_{p-2}^{(1)} \omega_{p-1}^{(0)}$$

and

$$x^{(j)} \to x^{(j+1)} \quad \text{when} \quad x^{(j)} \in X_i^{(j)}, \quad j < p-1,$$

$$x^{(p-1)} \to \delta_i(x)^{(0)} y \quad \text{when} \quad x^{(p-1)} \in X_i^{(p-1)},$$

$$y \to y.$$

Disregarding the non-commutativity of letters we may write

$$\begin{aligned} \omega &\to \delta_0(\omega_0)^{(0)} y^{\lg(\omega_0)} \omega_1^{(p-1)} \ldots \omega_{p-2}^{(2)} \omega_{p-1}^{(1)} \\ &\to \delta_0(\omega_0)^{(1)} y^{lg(\omega_0)} \delta_1(\omega_1)^{(0)} y^{\lg(\omega_1)} \ldots \omega_{p-1}^{(2)} \\ &\to \ldots \end{aligned}$$

Thus

$$\begin{aligned} r(n) &= (r(0)+s(0)+\ldots+s(m-1)) \\ &\quad +(s(m)+\ldots+s(m+p-1)) \\ &\quad +s(m+p)+\ldots+s(n-1) \\ &= (r(0)+s(0)+\ldots+s(m-1))+f_G(n-m-p), \end{aligned}$$

when $n \geqq m+p$. It is now easy to extend G to a PDOL system G' for which $f_{G'}(n) = r(n)$.

LEMMA 2: *Let $(r(n))$ be a Z-rational sequence. Then for any large natural number R the sequence defined by*

$$d(0) = d(1) = 1,$$

$$d(2n) = R^{2n} - r(2n),$$

$$d(2n+1) = R^{2n} - r(2n+1) \qquad (n > 0)$$

is a DOL *sequence.*

Proof: Let at first $(r(n))$ be a DOL sequence corresponding to the system $G = (X, \delta, \omega)$. Construct the system $H = (X \cup \bar{X} \cup \bar{\bar{X}} \cup \{a, b\}, \delta', \Omega)$ where

$$\Omega = a^{R^2-2r(2)-r(3)} \delta^2(\omega) \overline{\delta^3(\omega)}$$

and

$$x \to b\bar{\bar{x}}, \qquad \bar{x} \to \lambda, \qquad a \to b, \qquad b \to a^{R^2},$$

$$\bar{\bar{x}} \to a^{R^2(1+\lg(\delta(x)))-2\lg(\delta^2(x))-\lg(\delta^3(x))}\delta^2(x)\overline{\delta^3(x)}.$$

It is immediately seen that $f_H(n) = d(n+2)$.

This implies our lemma because of the following. It is known that every Z-rational sequence is the difference of two N-rational sequences (*see* [9] remark 2 or [2] p. 218). Furthermore, every N-rational sequence is the difference of two DOL sequences for

$$PM^n Q = (P+(1, \ldots, 1))M^n Q-(1, \ldots, 1)M^n Q.$$

Hence every Z-rational sequence can be written as the difference of two DOL sequences.

THEOREM 3: *Not every increasing* DOL *sequence is a* PDOL *sequence.*

Proof: Using lemma 2 we see that when R is a large natural number then the sequence $(d(n))$ where

$$d(0) = d(1) = 1,$$

$$d(2n) = R^{2n},$$

$$d(2n+1) = R^{2n}+(\mathrm{Re}(3+4i)^{2n+1})^2 \qquad (n > 0)$$

is an increasing DOL sequence. Now

$$\mathrm{Re}(3+4i)^n = \cos 2\pi n\alpha . 5^n,$$

where α is irrational because for every positive n $\mathrm{Im}(3+4i)^n < 4 \pmod 5$. The theorem of Berstel [1] then implies that the sequence $(d(2n+1)-d(2n))$ cannot be N-rational. Therefore $(d(n))$ is not a PDOL sequence.

Note: Let $(d(n)) = (PM^n Q)$ be a DOL sequence. By lemma 4 below we have $M^{m+p} \geqq M^m$ for some integers m and p $(p > 0)$. But then each of the sequences

$$(d(m+i+(n+1)p)-d(m+i+np)) \qquad (i = 0, \ldots, p-1)$$

is a DOL sequence and so the sequences

$$(d(m+i+np)) \qquad (i = 0, \ldots, p-1)$$

are PDOL sequences. This result also appears in [5] (proof of th. 4.12).

4. THE DOL SEQUENCES

Let $G = (\{x_1, \ldots, x_k\}, \delta)$ be a DOL scheme such that for any letter x_i:

$$\lg(\delta^n(x_i)) \sim g_i A^n \quad \text{as } n \to \infty \qquad (g_i > 0,\ A \geqq 1).$$

If $w \in \{ x_1, \ldots, x_k \}^+$ then the number

$$g(w) = \lg_1(w) g_1 + \ldots + \lg_k(w) g_k$$

is called the growth coefficient of w.

Suppose we have p DOL schemes $G_i = (X_i, \delta_i)$ $(i = 0, \ldots, p-1)$ satisfying the condition of the above definition with a common number A. Introduce the infinite alphabet

$$X = \{ (W_0, \ldots, W_{p-1}) \mid W_i \in X_i^+ \}$$

and define

$$\delta(W_0, \ldots, W_{p-1}) = (\delta_0 W_0, \ldots, \delta_{p-1} W_{p-1})$$

Define further

$$\pi(W_0, \ldots, W_{p-1}) = (\eta(W_0), \ldots, \eta(W_{p-1})),$$

where η means Parikh-vector. Take then a fixed element $(\omega_0, \ldots, \omega_{p-1})$ of X and denote

$$Y = \left\{ (W_0, \ldots, W_{p-1}) \,\middle|\, \frac{g(W_0)}{g(\omega_0)} = \ldots = \frac{g(W_{p-1})}{g(\omega_{p-1})} \right\}.$$

Obviously $(W_0, \ldots, W_{p-1}) \in Y$ implies that also $\delta(W_0, \ldots, W_{p-1}) \in Y$.

LEMMA 3 : *There are vectors $V_1, \ldots, V_J$ of $\pi(Y)$ such that any vector in $\pi(Y)$ is a sum of these.*

Proof: Let $\pi(Y) \subseteqq N^k$. By the definition of Y there are algebraic numbers a_{ij} $(i = 1, \ldots, p-1; j = 1, \ldots, k)$ such that $\bar{v} = (n_1, \ldots, n_k) \in Z^k$ is in $\pi(Y)$ iff

$$\bar{v} \neq 0, \qquad \bar{v} \geqq 0,$$

and

$$a_{i1} n_1 + \ldots + a_{ik} n_k = 0 \qquad (i = 1, \ldots, p-1).$$

Let F be the additive subgroup of Z^k defined by the above linear system.

We shall need the following simple lemma (*see* [3]):

LEMMA 4: *Any subset of N^k contains only a finite number of minimal vectors (with respect to the natural componentwise ordering).*

Let now $\bar{V}_1, \ldots, \bar{V}_J$ be the minimal vectors of $\pi(Y)$. If $\bar{V} \in \pi(Y)$ then it has a representation $\bar{V} = \bar{V}_i + \bar{U}$ where $\bar{U} \in N^k$. But if $\bar{U} \neq \bar{0}$ it is in $\pi(Y)$ because it belongs to F. Repeating this process we obtain $\bar{V}$ as a sum of the minimal vectors.

THEOREM 4: *Let $G_i = (X_i, \delta_i, \omega_i)$ $(i = 0, \ldots, p-1)$ be* DOL *systems such that if $x_j \in X_i$ then*

$$\lg(\delta_i^n(x_j)) \sim g_{ij} A^n \quad \text{as } n \to \infty \qquad (g_{ij} > 0,\ A \geqq 1).$$

Then the sequence defined by

$$d(np+i) = \lg(\delta_i^n(\omega_i)) \qquad (n = 0, 1, \ldots, i = 0, \ldots, p-1)$$

is a DOL *sequence.*

Proof: Take p copies $X_i^{(0)}, \ldots, X_i^{(p-1)}$ of each X_i and define

$$Y = \bigcup_{j=0}^{p-1} \left\{ (W_0^{(j)}, \ldots, W_{p-1}^{(j)}) \mid W_i^{(j)} \in X_i^{(j)+}, \quad \frac{g(W_0)}{g(\omega_0)} = \ldots = \frac{g(W_{p-1})}{g(\omega_{p-1})} \right\}.$$

Let $V_1^{(0)}, \ldots, V_J^{(0)}, \ldots, V_1^{(p-1)}, \ldots, V_J^{(p-1)}$ be elements of Y corresponding to the vectors given by lemma 3. We may say that any element of Y is a commutative product of these elements.

Introduce the DOL system

$$G = (\{ V_1^{(0)}, \ldots, V_J^{(p-1)} \} \cup \{ y \}, \delta, \omega) = (Z \cup \{ y \}, \delta, \omega),$$

where δ and ω are as follows:

ω consists of $(\omega_0^{(0)}, \ldots, \omega_{p-1}^{(0)})$ written commutatively in the alphabet Z and of so many y's that $\lg(\omega)$ becomes equal to $\lg(\omega_0)$;

y produces λ;

when $j < p-1$ $(W_0^{(j)}, \ldots, W_{p-1}^{(j)})$ produces $(W_0^{j+1}, \ldots, W_{p-1}^{(j+1)}$ and so many y's that the length of the produced word will be $\lg(W_{j+1})$;

$(W_0^{(p-1)}, \ldots, W_{p-1}^{(p-1)})$ produces $(\delta_0(W_0)^{(0)}, \ldots, \delta_{p-1}(W_{p-1})^{(0)})$ written commutatively in the alphabet Z and so many y's that the produced word will have length $\lg(\delta_0(W_0))$.

Heuristically, derivations in the systems G_i are simulated in the components of the letters of Z. With the aid of the p copies taken of the alphabets the simulation is delayed to happen only at intervals of p steps. By using the letter y the length of the word $\delta^{i+np}(\omega)$ is adjusted to be equal to that of $\delta_i^n(\omega_i)$. This is possible because the components of the letters of Z are nonempty words.

It should be clear now that $(d(n))$ is the growth sequence of G.

Let $G = (X, \delta)$ be a DOL scheme which gives a growth of order $n^l A^n$ $(A \geqq 1, l \geqq 0)$ but does not give a growth of higher order. We divide X into classes $\Sigma, \Sigma_0, \ldots, \Sigma_l$ as follows: the letters of Σ generate a growth having smaller order than A^n and the letters of Σ_i generate a growth of order $n^i A^n$.

It is clear that a letter of Σ_i cannot produce letters of $\Sigma_{i+1} \cup \ldots \cup \Sigma_l$, it must produce a letter of Σ_i and it may produce letters of $\Sigma \cup \Sigma_0 \cup \ldots \cup \Sigma_{i-1}$; a letter of Σ may produce only letters of Σ or λ.

LEMMA 5: *Any letter of Σ_l $(l > 0)$ generates letters of Σ_{l-1}. If all letters of $\Sigma \cup \Sigma_0 \cup \ldots \cup \Sigma_{l-1}$ are deleted then the resulting scheme H is such that all letters generate a growth of order A^n.*

Proof: Suppose x generates in H a growth whose order is at least $n A^n$. Then x generates in G in $2n$ steps a word whose length is at least of the order

$$n A^n n^l A^n = \left(\frac{1}{2}\right)^{l+1} (2n)^{l+1} A^{2n}.$$

This shows that the growth of x in H has the order A^n at most.

Assume that $x \in \Sigma_l$ $(l > 0)$ never generates a letter of Σ_{l-1}. By the above $\delta^n(x)$ contains $O(A^n)$ letters of $\Sigma \cup \Sigma_0 \cup \ldots \cup \Sigma_{l-2}$ directly produced by letters of Σ_l. But

$$\sum_{n=0}^{N} A^n (N-n)^{l-2} A^{N-n} = A^N \sum_{n=0}^{N} n^{l-2} = o(N^l A^N)$$

and so x cannot generate a growth of order $n^l A^n$. Hence x must generate letters of Σ_{l-1}.

Assume further that the growth of x in H is majorized by a^n $(a < A)$. Because

$$\sum_{n=0}^{N} a^n (N-n)^{l-1} A^{N-n} = a^N \sum_{n=0}^{N} n^{l-1} (A/a)^n = O(N^{l-1} A^N)$$

we have a contradiction as above. Thus the growth of x in H has the order A^n when $l > 0$.

Suppose $x \in \Sigma_0$ and the growth of x in H as well as the growth of any letter of Σ in G is majorized by a^n $(a < A)$. Because

$$\sum_{n=0}^{N} a^n a^{N-n} = o(A^N)$$

we see that the above result is true also when $l = 0$.

THEOREM 5: *Let $(r(n))$ be an N-rational sequence such that $r(n) \neq 0$ for every n and the quotient $r(n+1)/r(n)$ remains bounded. Then $(r(n))$ is a* DOL *sequence.*

Proof: We know that there are numbers m and p and DOL sequences $(d_0(n)), \ldots, (d_{p-1}(n))$ such that

$$r(m+i+np) = d_i(n).$$

By our assumption all these sequences have the same order of growth. We may suppose that it is of the form $n^l A^n$ $(A > 1)$ for the polynomial case is covered by a theorem of Ruohonen [7].

Let $G_i = (X_i, \delta_i, \omega_i)$ be a DOL system corresponding to the sequence $(d_i(n))$. Write $X_i = \Sigma_i \cup \Sigma_{i0} \cup \ldots \cup \Sigma_{il}$ as before and denote by H_{ij} the DOL schema obtained from (X_i, δ_i) by deleting all letters except those of Σ_{ij}.

Suppose $W \in \Sigma_{i0}^{+}$ and $w \in \Sigma_i^{*}$. If W generates Vv ($V \in \Sigma_{i0}^{+}$, $v \in \Sigma_i^{*}$) and w generates u in k steps then there are constants N, M and L (independent of k and i) such that

$$\begin{aligned} \lg(V) &\geqq NA^k \lg(W), \\ \lg(v) &\leqq MA^k \lg(W), \\ \lg(u) &\leqq La^k \lg(w) \qquad (a < A). \end{aligned}$$

We now see that if $\lg(w)/\lg(W) \leqq \alpha\, M/N$ ($\alpha \geqq 2$, $\alpha M/N$ integer) and if k is so large that $(L/N)\ (a/A)^k \leqq 1/2$ then

$$\begin{aligned} \lg(uv)/\lg(V) &\leqq M/N + (L/N)(a/A)^k (\lg(w)/\lg(W)) \\ &\leqq M/N + (1 - 1/\alpha)\alpha\, M/N = \alpha\, M/N, \end{aligned} \tag{1}$$

too.

By taking a multiple of p, if necessary, we may suppose that the following three conditions hold:

(i) any growth in H_{ij} ($i = 0, \ldots, p-1$; $j = 0, \ldots, l$) is asymptotically equal to constant times A^n;

(ii) any letter of Σ_{ij} ($i = 0, \ldots, p-1$; $j = 0, \ldots, l$) produces in a step letters of all the alphabets $\Sigma_{i,j-1}, \ldots, \Sigma_{i,0}$;

(iii) the equation (1) holds with $k = 1$.

Moreover, by increasing m we obtain the following situation:

(iv) the axiom of G_i contains letters of all the alphabets $\Sigma_{i,0}, \ldots, \Sigma_{il}$.

We are now ready to give an induction proof showing that the sequence $(s(n)) = (r(n-m))$ is a DOL sequence. This immediately implies our theorem.

If $l = 0$ we at first neglect all letters of the Σ_i's and construct a system just as in the proof of theorem 4. Then we take p copies $x^{(0)}, \ldots, x^{(p-1)}$ of each neglected letter x and join these copies to the components of the letters of Z so that the original systems G_i become simulated. Condition (iii) assures that this can be done. At the same time we add y's so that the right lengths are obtained.

When taking the induction step we at first delete all letters of the alphabets $\Sigma_i \cup \Sigma_{i0} \cup \ldots \cup \Sigma_{i,l-1}$ and construct a system according to theorem 4. By conditions (ii) and (iv) the letters of this system as well as the p-tuple $(\omega_0, \ldots, \omega_{p-1})$ give axioms for systems whose existence is guaranteed by the induction hypothesis.

Example: Let

$$G_0 = (\{A, B, C\}, \delta_0, AB),$$

where

$$A \to A^4 B, \qquad B \to B^4 b, \qquad b \to b$$

and

$$G_1 = (\{C, D, E, F\}, \delta_1, CD),$$

where

$$C \to C^4 D, \qquad D \to D^2 E^3, \qquad E \to E^2 F^4, \qquad F \to DF.$$

The common order of growth is $n.4^n$ and

$$\Sigma_0 = \{b\}, \qquad \Sigma_{00} = \{B\}, \qquad \Sigma_{01} = \{A\},$$
$$\Sigma_1 = \varnothing, \qquad \Sigma_{10} = \{D, E, F\}, \qquad \Sigma_{11} = \{C\}.$$

The procedures described in the preceding proof yield e. g. the following system: The axiom is

$$(A^{(0)}, C^{(0)})(B^{(0)}, D^{(0)})$$

and the productions are

$$(A^{(0)}, C^{(0)}) \to (A^{(1)}, C^{(1)}),$$
$$(A^{(1)}, C^{(1)}) \to (A^{(0)}, C^{(0)})^4 (B^{(0)}, D^{(0)}),$$
$$(B^{(0)}, D^{(0)}) \to (B^{(1)}, D^{(1)}),$$
$$(B^{(1)}, D^{(1)}) \to (B^{(0)} b^{(0)}, D^{(0)})(B^{(0)}, D^{(0)})(B^{(0)2}, E^{(0)3}) y^2.$$
$$(B^{(0)2}, E^{(0)3}) \to (B^{(1)2}, E^{(1)3}) y^2,$$
$$(B^{(1)2}, E^{(1)3}) \to (B^{(0)2} b^{(0)}, E^{(0)3})^2 (B^{(0)}, F^{(0)3})^4 y^4,$$
$$(B^{(0)}, F^{(0)3}) \to (B^{(1)}, F^{(1)3}) y^2,$$
$$(B^{(1)}, F^{(1)3}) \to (B^{(0)} b^{(0)}, F^{(0)3})(B^{(0)}, D^{(0)})^3 y,$$
$$(B^{(0)} b^{(0)}, D^{(0)}) \to (B^{(1)} b^{(1)}, D^{(1)}),$$
$$(B^{(1)} b^{(1)}, D^{(1)}) \to (B^{(0)} b^{(0)}, D^{(0)})^2 (B^{(0)2}, E^{(0)3}) y^3,$$
$$(B^{(0)2} b^{(0)}, E^{(0)3}) \to (B^{(1)2} b^{(1)}, E^{(1)3}) y^2,$$
$$(B^{(1)2} b^{(1)}, E^{(1)3}) \to (B^{(0)2} b^{(0)}, E^{(0)3})^2 (B^{(0)} b^{(0)}, F^{(0)3})(B^{(0)}, F^{(0)3})^3 y^5,$$
$$(B^{(0)} b^{(0)}, F^{(0)3}) \to (B^{(1)} b^{(1)}, F^{(1)3}) y^2,$$
$$(B^{(1)} b^{(1)}, F^{(1)3}) \to (B^{(0)} b^{(0)}, F^{(0)3})(B^{(0)} b^{(0)}, D^{(0)})(B^{(0)}, D^{(0)})^2 y^2,$$
$$y \to \lambda.$$

Note: Let $(d(n))$ be a DOL sequence such that the rational function $\sum d(n) x_n$ is not a polynomial. Then it is easy to see that the growth order of $(d(n))$ is $n^l A^n$ where $1/A$ is the smallest positive pole of $\sum d(n) x^n$ and $l+1$

is its multiplicity. This enables us to effectively compare the growth orders of two DOL sequences; we describe the method briefly in general form.

Given integer polynomials $q_1(x), \ldots, q_p(x)$ we can (using symmetric polynomials) construct a polynomial $Q(x)$ such that any difference of two zeros of $q(x) = q_1(x) \ldots q_p(x)$ is a zero of $Q(x)$. Thus we can give a positive number γ such that if z_1 and z_2 are zeros of $q(x)$ then either $z_1 = z_2$ or $|z_1 - z_2| > \gamma$.

The polynomial

$$Q_i(x) = \frac{q_i(x)}{g.c.d.(q_i(x), q_i'(x))}$$

has simple zeros which are the same as those of $q_i(x)$. Therefore we can compare the real roots of the polynomials $q_i(x)$ by examining the sign changes of the polynomials $Q_i(x)$.

Because α is a k-fold zero of $q_i(x)$ iff it is a zero of $q_i(x), q_i'(x), \ldots, q_i^{(k-1)}(x)$ but not a zero of $q^{(k)}(x)$ it is possible to determine the multiplicity of any real zero of $q_i(x)$.

An N-rational sequence $(r(n))$ is by theorem 5 a DOL sequence iff one of the following conditions holds:

(i) there is a natural number L such that $r(n) > 0$ when $n \leqq L$ and $r(n) = 0$ when $n > L$;

(ii) every $r(n)$ is positive and the DOL sequences given by theorem 1 have the same growth order.

Hence it is possible to decide whether or not a given N-rational sequence is a DOL sequence.

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